

Two Birds, One Stone: A Fast, yet Lightweight, Indexing Scheme for Modern Database Systems

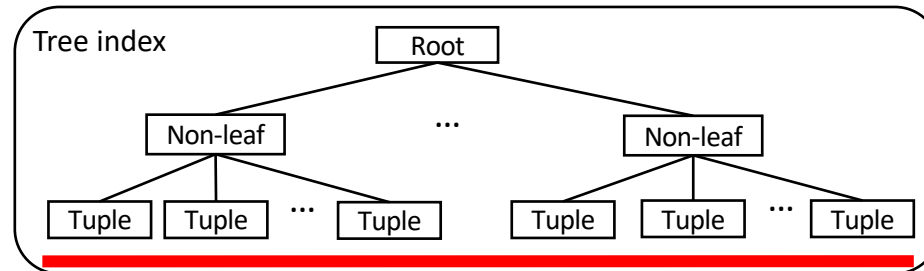
Jia Yu and Mohamed Sarwat

Arizona State University

Motivation

- **Tree index**

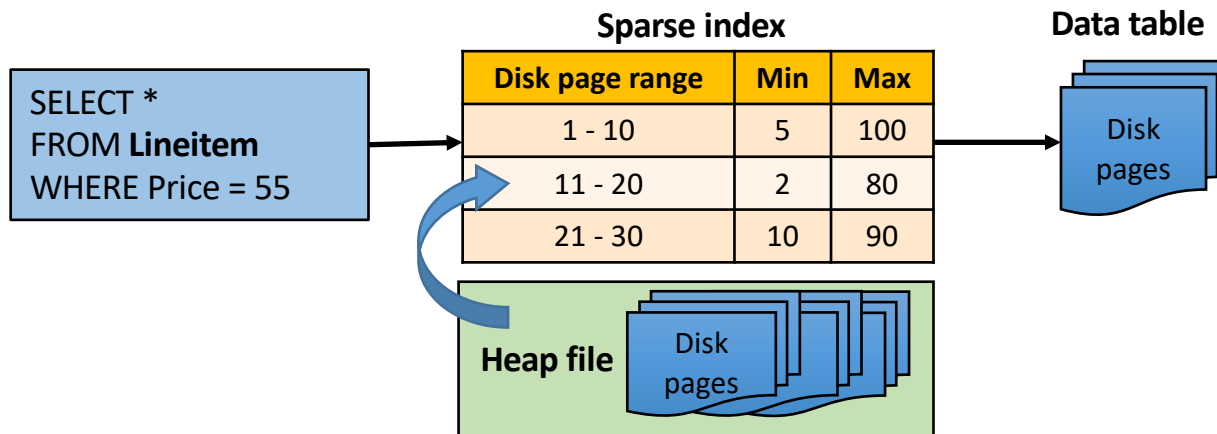
- Fast index search
- Large storage overhead: 5% - 15% additional overhead
- Slow maintenance: Split, merge, redistribute nodes



B+ Tree on TPC-H	Index size	Initialization time	Insertion time (0.1%)
2 GB	0.25 GB	30 sec	10 sec
20 GB	2.51 GB	500 sec	1180 sec
200 GB	25 GB	8000 sec	42000 sec

Motivation (cont.)

- **Bitmap index:** Fast, small, low cardinality, read-only
- **Compressed index:** Small, slow search, slow maintenance
- **Approximate index:** Inaccurate result
- **Sparse index** (aggregated page info: min, max): Very small, not good at queries



Design goals

- **Low indexing overhead** (size and initialization time)
- **Fast index maintenance** (insertion and deletion)
- **Competitive query performance**
 - No need to be fast on all selectivity scenarios
 - Works for “**not too selective**” but still “**selective**” query

TPC-H Lineitem Table

ID	QUANTITY	PRICE	SHIPDATE	Comment
1	4	\$50	1995-08-01	N/A
2	22	\$90	1994-11-27	Refurbished
...	...	...	...	...

Highly selective, <0.001%

```
SELECT *  
FROM Lineitem  
WHERE Price > 0.00 AND Price < 0.01
```

Use B+ Tree

Still selective 0.01% - 1%

```
SELECT average(*)  
FROM Lineitem  
WHERE Price > 55 AND Price < 65
```

Use our index

Unselective > 1%

```
SELECT *  
FROM Lineitem  
WHERE Price > 10
```

Full table scan

Hippo Index Structure

- Sparse index

Complete load balanced histogram

Bucket	Price
1	1 - 20
2	21 - 40
3	41 - 60
4	61 - 90
5	91 - 120

Manage



Hippo logical structure		
Disk Page #	Partial histogram	Internal data
1 - 10	2,3,4	21,22,55,75,77
11 - 25	2,4,5	23,24,62,91,92
26 - 30	1,2,5	11,12,25,101,110

Grouped by partial histogram similarity

Bucket	Price
2	21 - 40
3	41 - 60
4	61 - 90

Partial histogram 1

Bucket	Price
2	21 - 40
4	61 - 90
5	91 - 120

Partial histogram 2

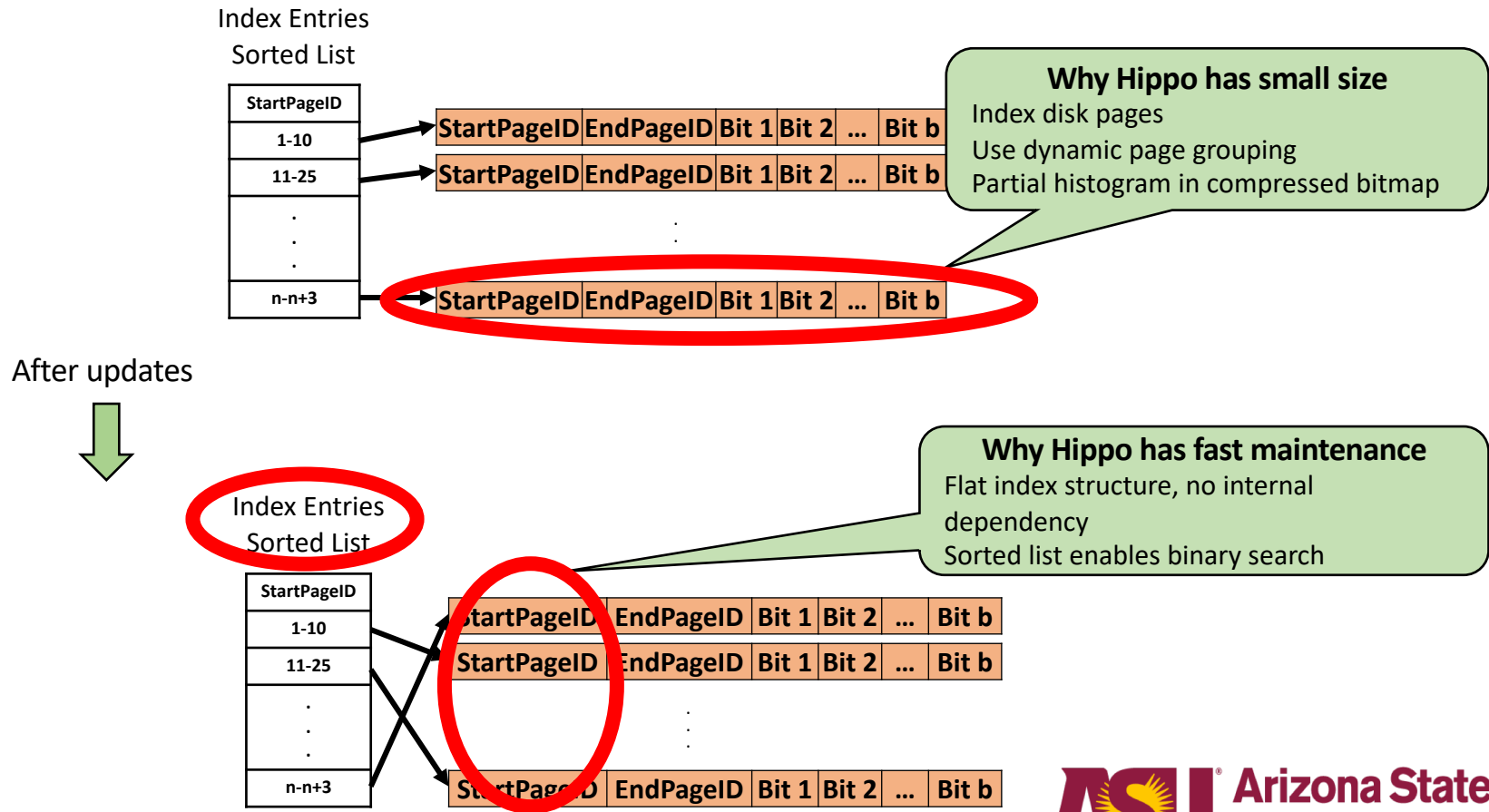
Bucket	Price
1	1 - 20
2	21 - 40
5	91 - 120

Partial histogram 3

- ✓ No initial. overhead
- ✓ Load-balanced, handle data skewness
- ✓ Global data distribution won't change frequently

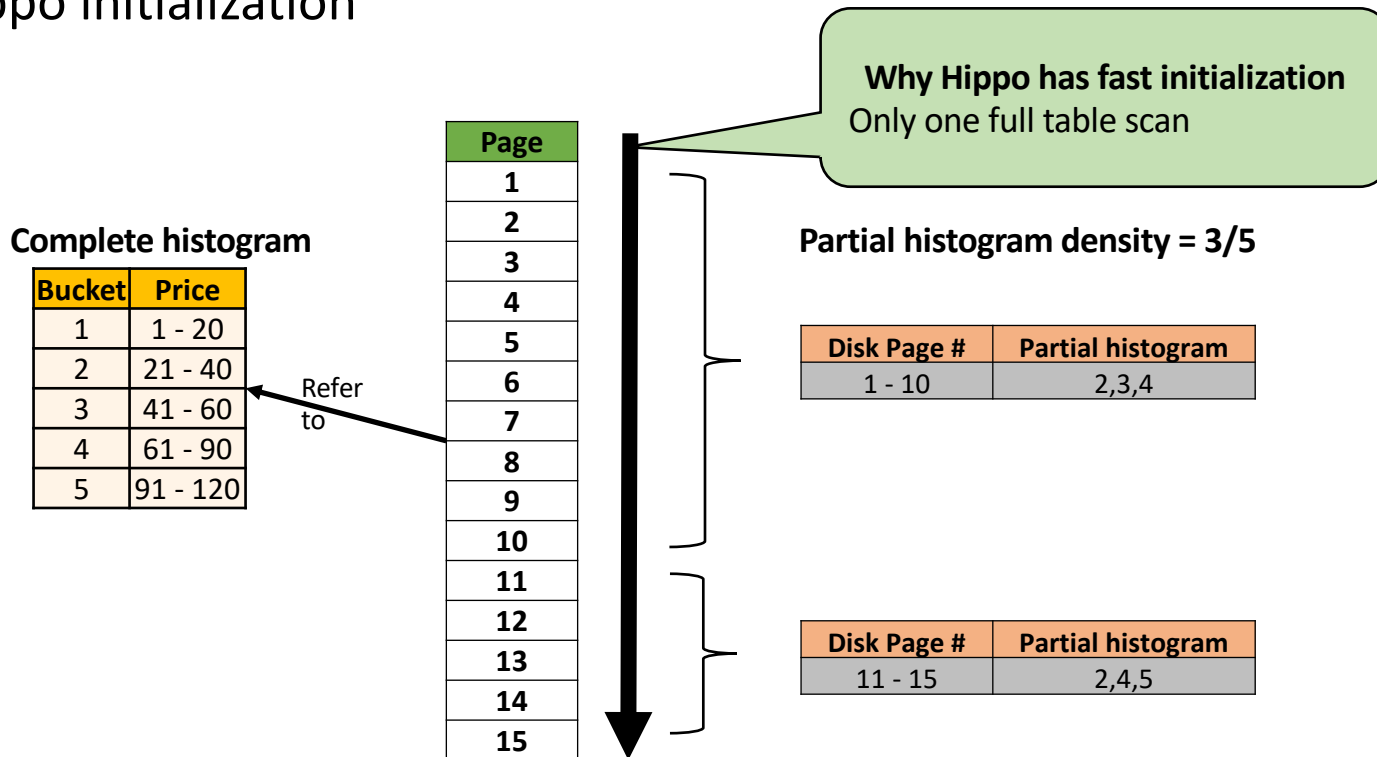
Hippo Index Structure (cont.)

- Hippo disk structure



Hippo Index Structure (cont.)

- Hippo initialization



Hippo Index Search

- Step 1: Scanning index entries

```
SELECT *  
FROM Lineitem  
WHERE Price = 55
```

Bucket	Price
1	1 - 20
2	21 - 40
3	41 - 60
4	61 - 90
5	91 - 120

```
SELECT *  
FROM Lineitem  
WHERE Price > 55
```

Bucket	Price
1	1 - 20
2	21 - 40
3	41 - 60
4	61 - 90
5	91 - 120

```
SELECT *  
FROM Lineitem  
WHERE Price > 55 AND Price < 65
```

Bucket	Price
1	1 - 20
2	21 - 40
3	41 - 60
4	61 - 90
5	91 - 120

Hippo Index Search

- Step 1: Scanning index entries (continued)

SELECT *
FROM Lineitem
WHERE Price = 55

Bucket	Price
1	1 - 20
2	21 - 40
3	41 - 60
4	61 - 90
5	91 - 120

Why Hippo has fast query performance
Use partial histogram to prune useless disk pages

Hippo	
Disk Page #	Partial histogram
1 - 10	2,3,4
11 - 25	2,4,5
26 - 30	1,2,5
...	

Page 1 - 10

Bucket	Price
2	21 - 40
3	41 - 60
4	61 - 90

Partial histogram 1

Bucket	Price
2	21 - 40
4	61 - 90
5	91 - 120

Partial histogram 2

Bucket	Price
1	1 - 20
2	21 - 40
5	91 - 120

Partial histogram 3

Hippo Index Search

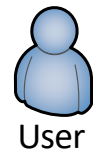
- Step 2: Filtering false positive pages

```
SELECT *  
FROM Lineitem  
WHERE Price = 55
```

Page 1 - 10

Page	Content	False positive
1	21,22,77	√
2	22,75,77	√
3	22,55,75	Got results!
4	21,55,75	Got results!
5	21,22,75	√
6	22,55,77	Got results!
7	21,22,77	√
8	21,55,77	Got results!
9	55,75,77	Got results!
10	21,75,77	√

Query
result



Hippo Maintenance: Insertion - Eager

- **Check index entries right after each data insertion**
 - Check the complete histogram and each Hippo partial histogram
- **Update index entries right after each check**
 - S1. No need to update the corresponding index entry
 - S2. Need to update the corresponding index entry at once

Hippo Maintenance: Insertion - Eager (cont.)

- Scenario 1: No need to update the entry

INSERT INTO Lineitem (Quantity, Price, Shipdate)
VALUES (3, **23**, 1997-05-30);

Why Hippo has fast maintenance
Sorted list enables binary search on index entries

Complete histogram

Bucket	Price
1	1 - 20
2	21 - 40
3	41 - 60
4	61 - 90
5	91 - 120

Page 1

Binary search

Hippo

Disk Page #	Partial histogram
1 - 10	2,3,4
11 - 25	2,4,5
26 - 30	1,2,5
...	

Bucket	Price
2	21 - 40

Disk Page #	Partial histogram
1 - 10	2 ,3,4

No index updates

Hippo Maintenance: Insertion - Eager (cont.)

- Scenario 2: Need to update the entry

```
INSERT INTO Lineitem (Quantity, Price, Shipdate)
VALUES (10, 61, 1995-01-01);
```

Complete histogram

Bucket	Price
1	1 - 20
2	21 - 40
3	41 - 60
4	61 - 90
5	91 - 120

Bucket	Price
4	61 - 90

Page
26

Binary search

Hippo

Disk Page #	Partial histogram
1 - 10	2,3,4
11 - 25	2,4,5

Why Hippo has fast maintenance

Flat index structure, no internal dependency

Only update one index entry

Update

Disk Page #	Partial histogram
26 - 30	1,2,5

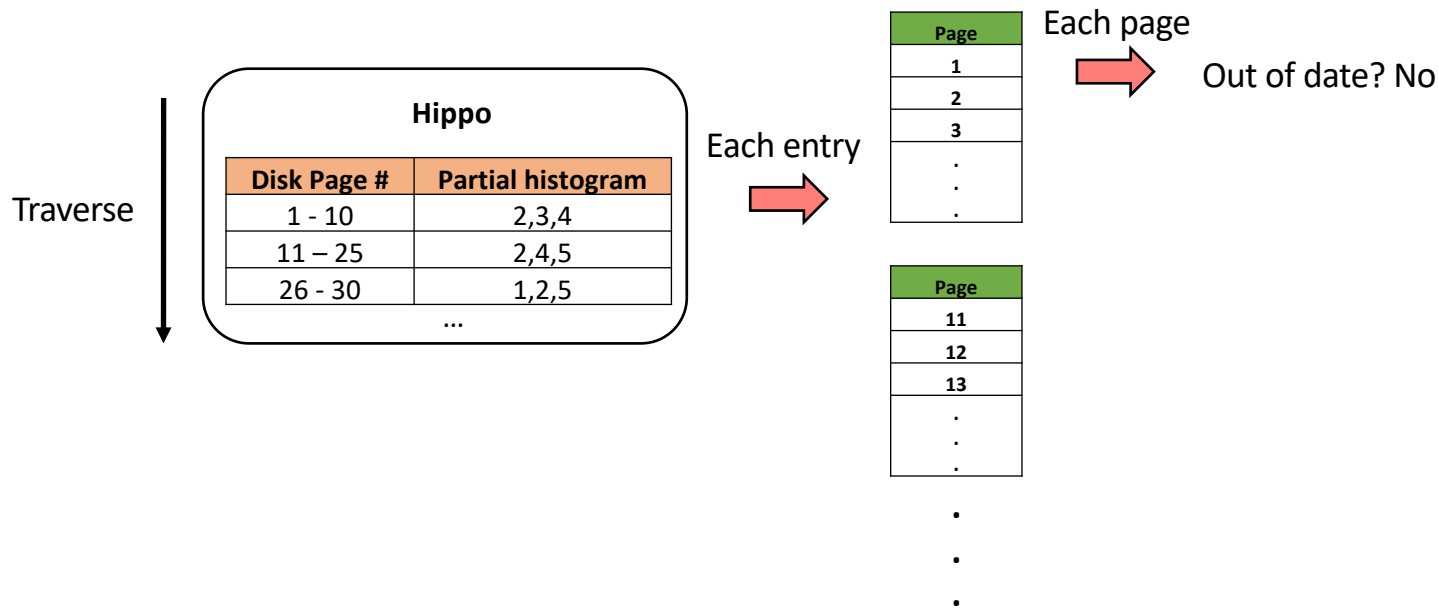
Disk Page #	Partial histogram
26 - 30	1,2, 4 ,5

Hippo Maintenance: Deletion - Lazy

- **Don't update index entries after each data deletion**
 - Deleted data tuples are marked as “deleted” and may be removed at any time
 - Still guarantee **accurate query results** because of “filtering false positive pages”
- **Update index periodically or invoked by the user**
 - Slow but can run at DBMS idle time
 - S1. No need to update the corresponding index entry
 - S2. Need to update the corresponding index entry

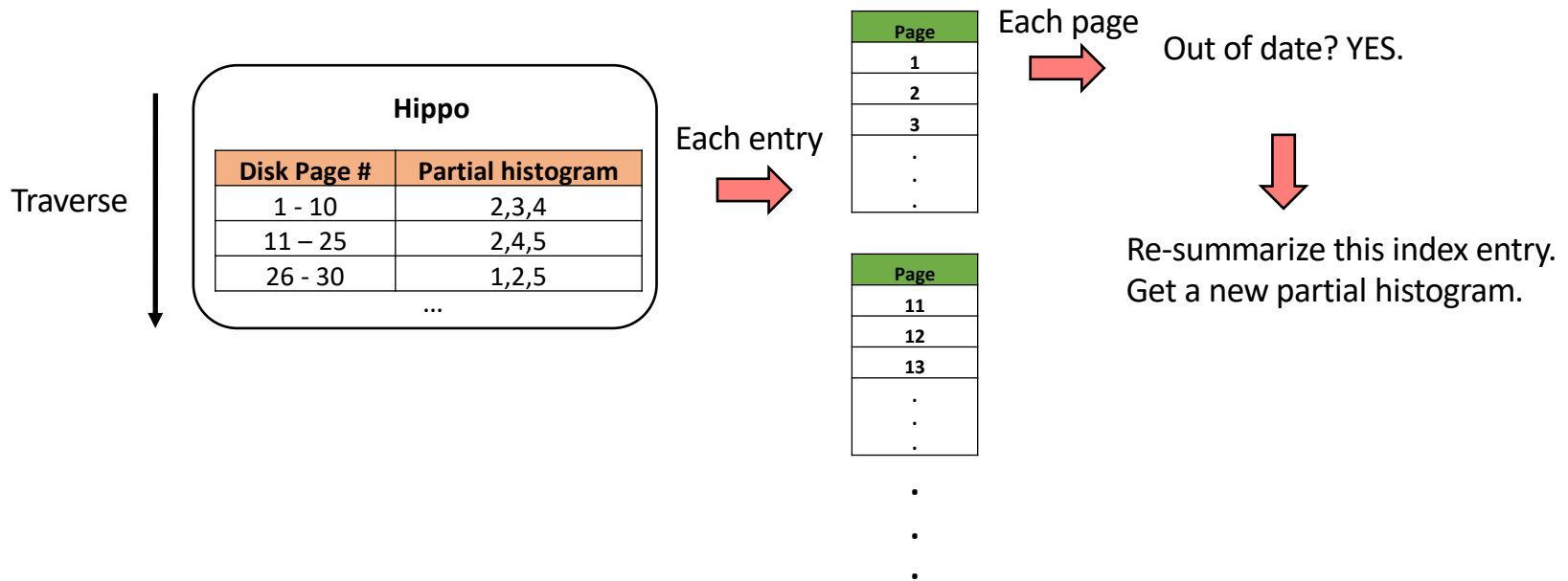
Hippo Maintenance: Deletion – Lazy (cont.)

- Scenario 1: No need to update the entry



Hippo Maintenance: Deletion – Lazy (cont.)

- Scenario 2: Need to update the entry



Experiments

- **Datasets (all around 200 GB):**

1. TPC-H: Lineitem table
 1. PartKey: Uniform distribution. 40 million distinct values
 2. SuppKey: Uniform distribution. 2 million distinct values
 3. OrderKey: sorted
2. Exponential distribution data: Skewed distribution
3. Wikipedia page hit rates: Skewed distribution
4. NYC Taxi trip records: pickup point, dimension reduction using Hilbert Curve

- **Compared indexes:**

1. B+ Tree
2. Hippo
3. [Block Range Index - BRIN \(A sparse index\)](#)

Experiments (cont.)

- Indexing overhead
 - Size: Hippo 40x < B+ Tree
 - Initial. time: Hippo 2.5x < B+ Tree

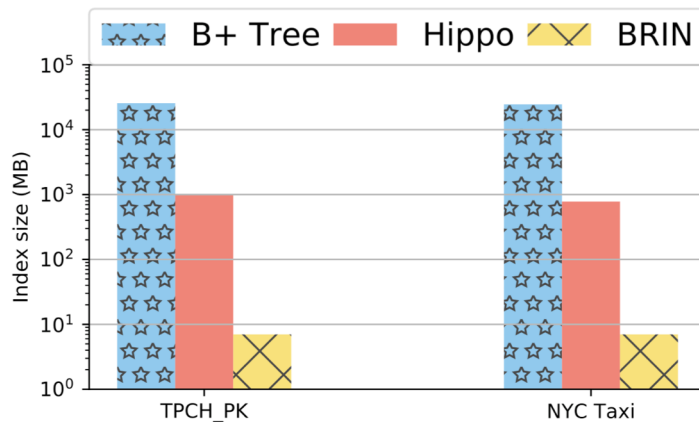


Figure: Index size on different datasets (logarithmic scale)

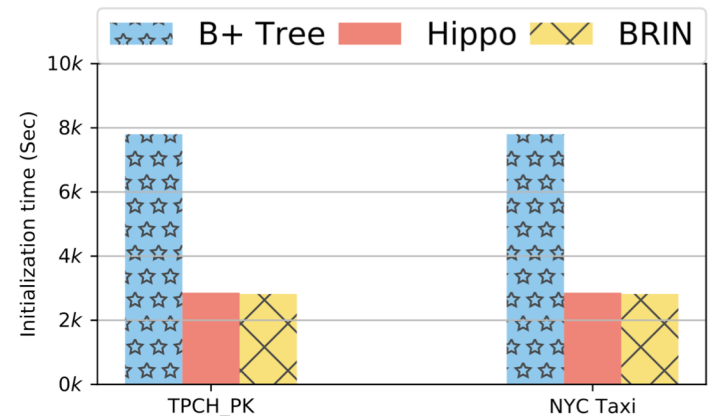


Figure: Initial. time on different datasets

Experiments (cont.)

- Query time
 - Hippo $\approx$ B+ Tree at 0.1% and 1% selectivity
 - BRIN always scans lots of disk pages

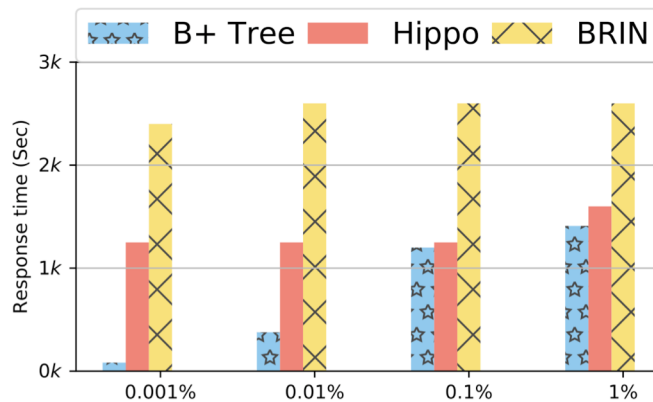


Figure: Query time on TPC-H

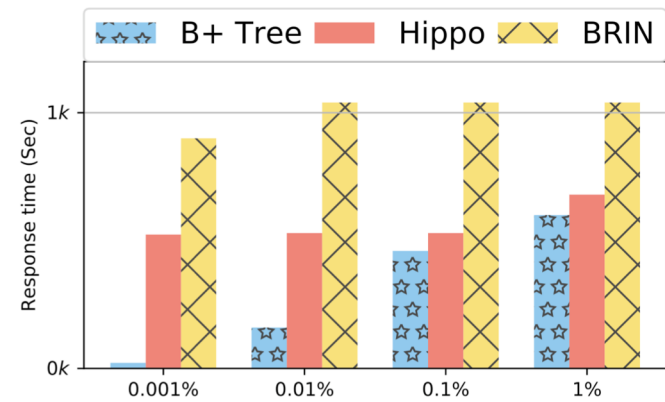


Figure: Query time on NYC Trips

Experiments (cont.)

- Hybrid workload
 - Update-intensive workloads (10%-50%), Hippo is 100x > B+ Tree
 - Query-intensive workloads (70%-90%), Hippo $\approx$ B+ Tree

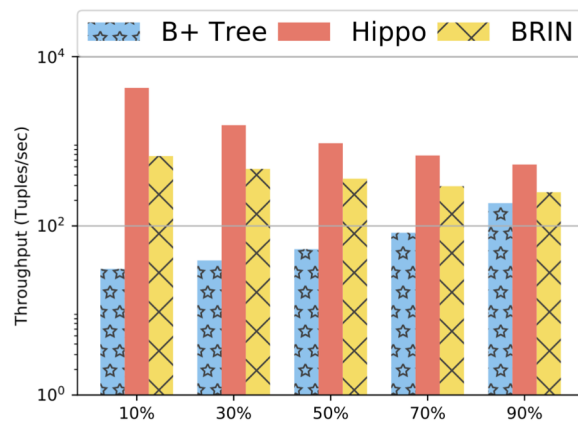


Figure: Throughput on TPC-H

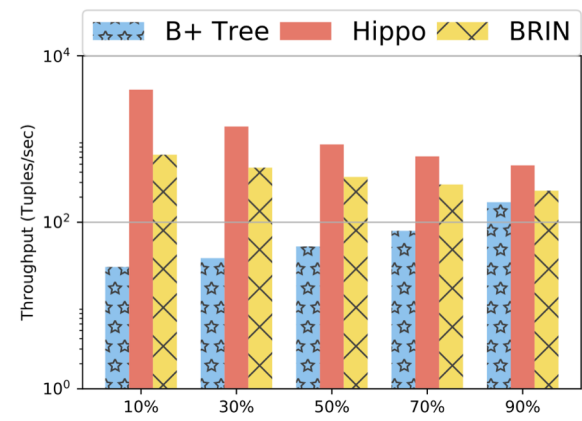


Figure: Throughput on NYC Trips

Take-home Lesson

- Use Hippo if you have limited storage budget
- Use Hippo if your query selectivity is 0.1% - 1%
- Use Hippo if you have update-intensive workloads

Questions?



PostgreSQL 9.6.1
kernel



<https://github.com/DataSystemsLab/hippo-postgresql>

Build Hippo

```
CREATE INDEX hippo_idx ON hippo_tbl USING hippo (randomNumber) WITH (density = 20);
```

Experiments (cont.)

- Maintenance time
 - Hippo 10x – 1000x < B+ Tree
 - Hippo 10x < BRIN. BRIN doesn't support deletion.

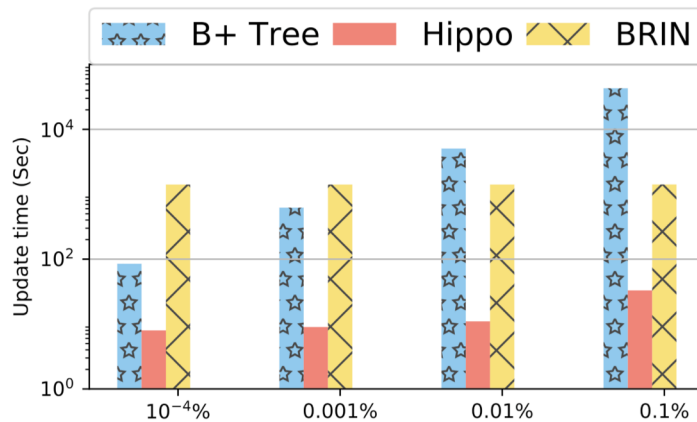


Figure: Update time on TPC-H
(logarithmic scale)

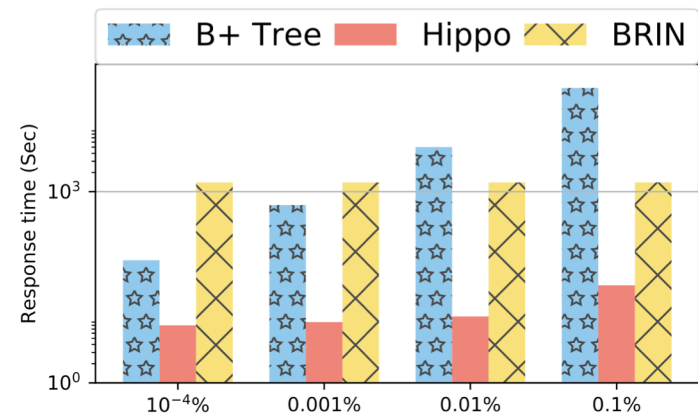


Figure: Update time on NYC Trips
(logarithmic scale)